

SwitchTx: Scalable In-Network Coordination for Distributed Transaction Processing Junru Li¹, Youyou Lu¹, Yiming Zhang², Qing Wang¹, Zhuo Cheng³, Keji Huang³, Jiwu Shu¹



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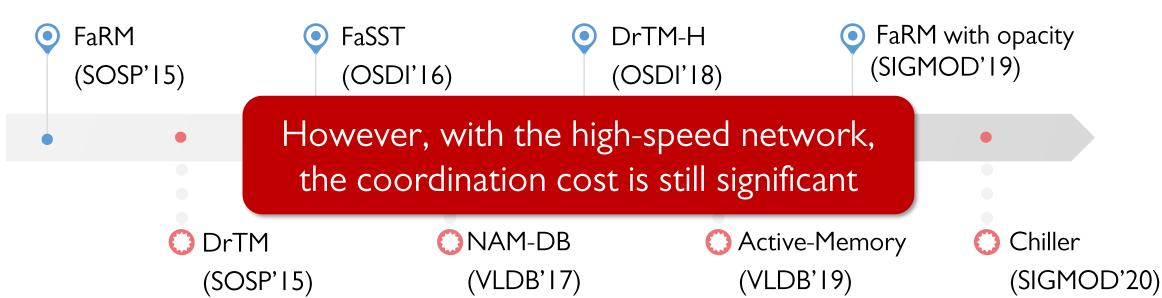
Coordination in Distributed Transactions

* Network communication is a major source of coordination cost

- Concurrency control protocols
- Replication protocols

* Leveraging the high-speed network

- Reduce latency
- Shorten contention span to reduce abort rate



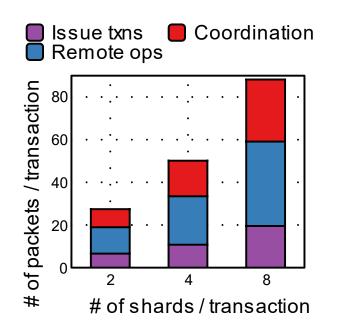
Coordination Cost I

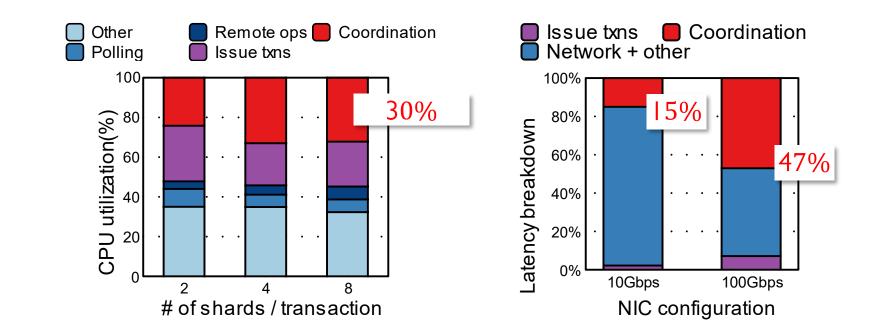
* With high-speed network, the coordination cost is still significant

Waste CPU to process coordination packets

Waste CPU cycles

CPU processing latency is more important with a faster network



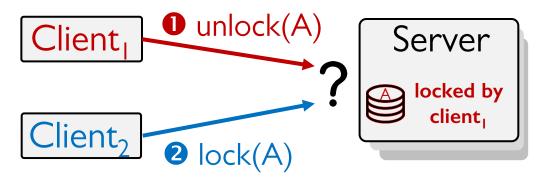


Coordination Cost II

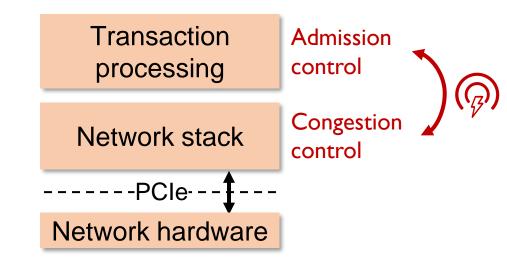
* With high-speed network, the coordination cost is still significant

Semantic gap between Txn apps and network

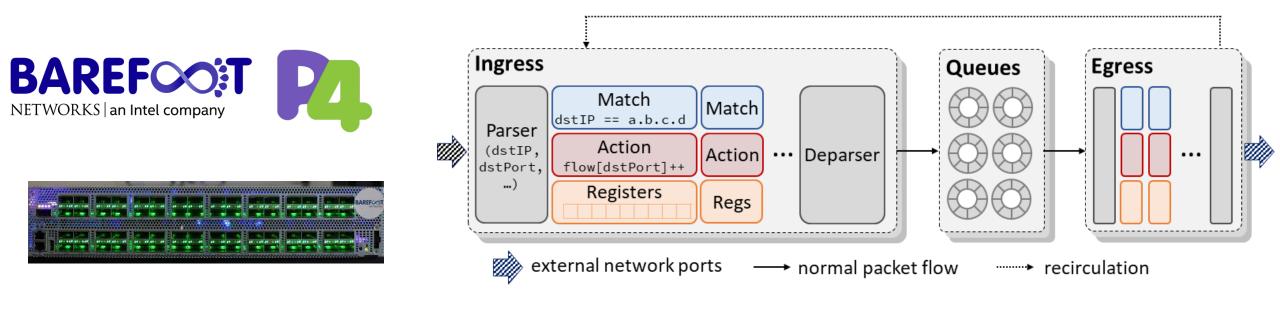
- Inappropriate processing order introduces extra aborts
- Redundant flow control algorithms interfere with each other
 Admission control: controls the number of concurrent transactions
 - Congestion control: controls the number of concurrent network messages



Situation **1**>2 : Client₂ locks A successfully Situation **2**>**1** : Client₂ needs to retry to lock A



Opportunities from Programmable Switches



Programmable Switches Centralized hub User-defined parsers / Match-Action tables / queues On-chip memory Line-rate processing

Design Goals and Challenges

Design Goals: reduce coordination cost

Offload coordination tasks

Manipulate transaction traffic intelligently

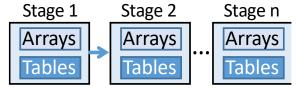
Challenges:

Restricted expressive power and limited on-chip memory

Coordination logic







Multi-switch scalability











Background & Motivation

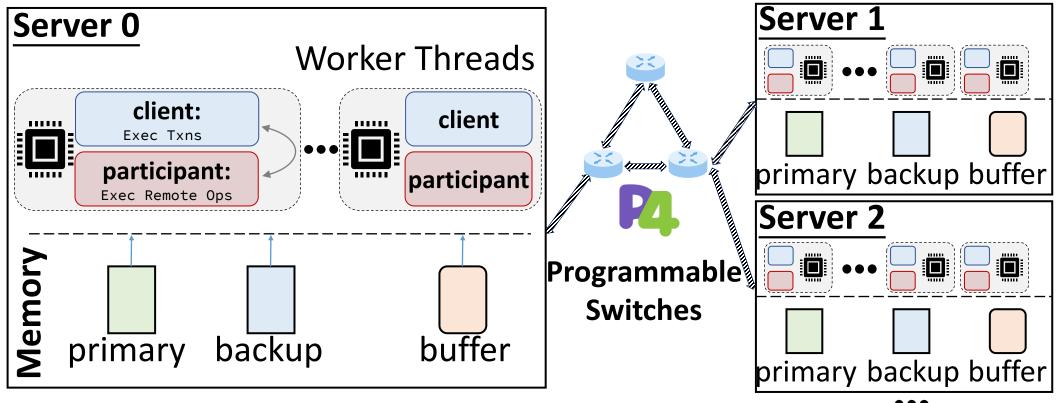
SwitchTx: In-Network Transaction Coordination

Results

Summary

Overview

In-Network Transaction Coordination



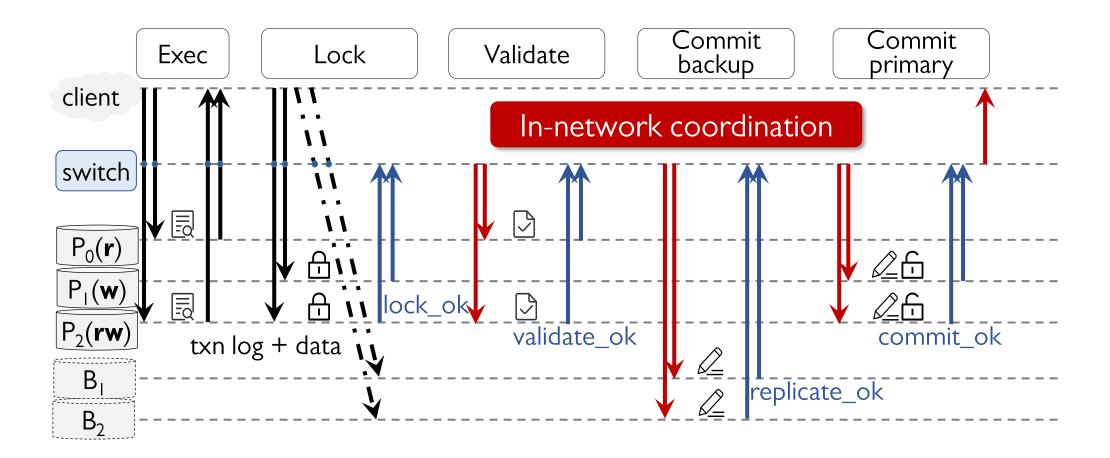


To save CPU utilization

- \bigcirc I.I Coordination tasks → in-switch Gather-and-Scatter (GaS)
- I.2 Scalable tree-based GaS using all switches
- To break the semantic gap between Txn apps and network
- 2.1 Semantic-aware packet priority control
- 2.2 Dynamic admission control

I.I) In-switch Gather-and-Scatter

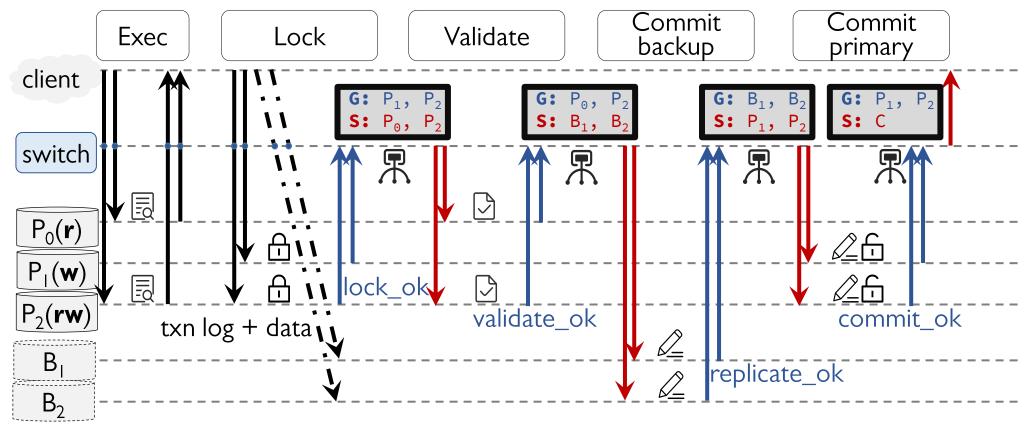
Offload coordination tasks as in-switch GaS
 An example: Txn { read[D₀,D₂], write[D₁,D₂] }



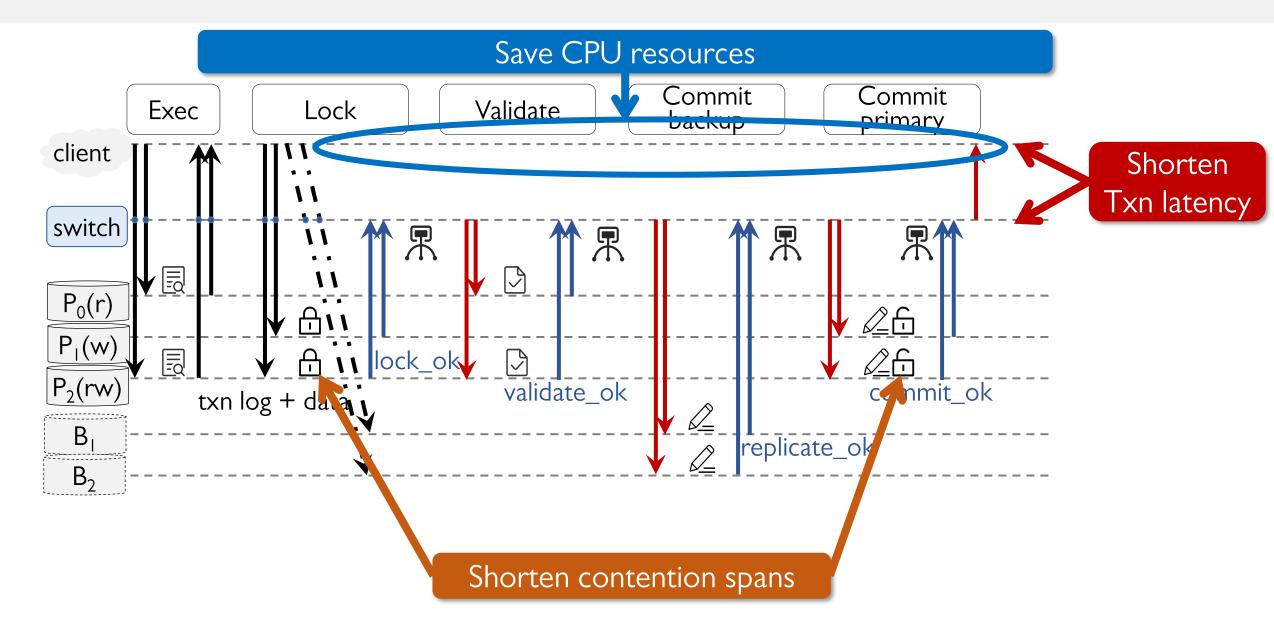
I.I) In-switch Gather-and-Scatter

* GaS (gather_group, scatter_group)

- * Gather messages from the participants of the current phase
- * Scatter the result to the participants of the next phase



I.I) In-switch Gather-and-Scatter

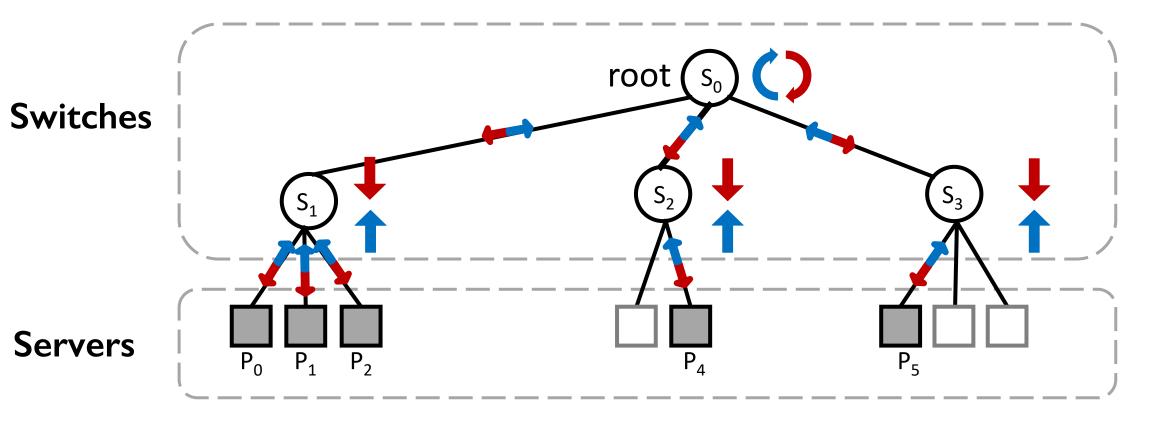


I.2) Scalable Tree-based Gather-and-Scatter

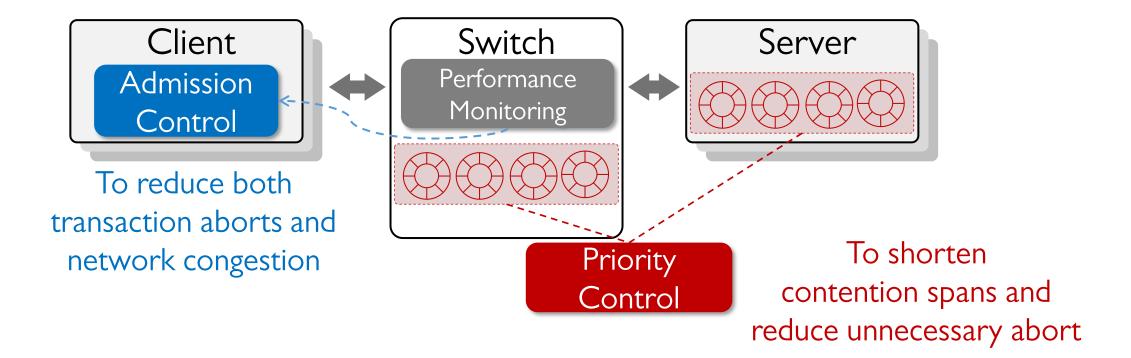
Gather-and-Scatter tree

- Servers: leaf nodes
- Switches: non-leaf nodes





2 Break the semantic gap

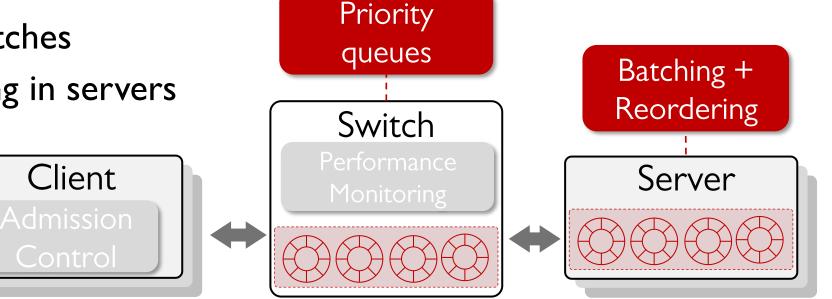


2.1) Semantic-aware Packet Priority Control

- * Assign priorities to messages based on their types
- Highest: lock releasing + messages of retrying transactions
- Lowest: lock acquiring
- Medium: other messages

Implementation

- Priority queues in switches
- Batch-based reordering in servers



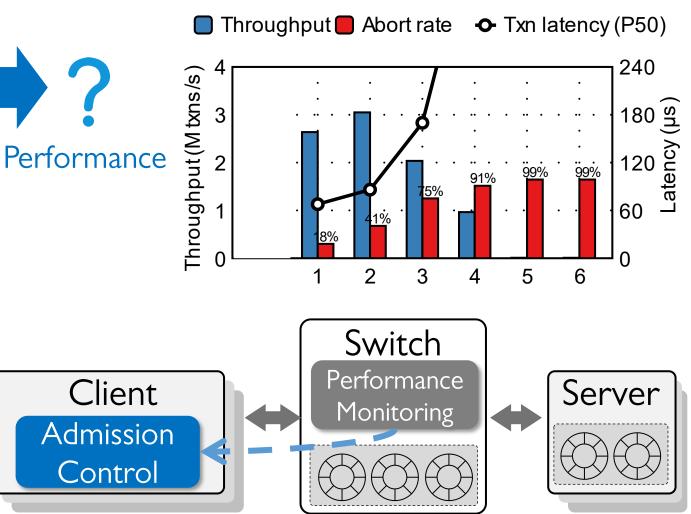
2.2) Dynamic Admission Control

Increasing maximum number of parallel requests

- ✤ Higher resource utilization
- Higher abort rate
- Network congestion

* Signals

- Global performance metrics
- Individual network conditions
- Algorithm: AIMD
 - Additive increase
 - Multiplicative decrease



More Details: checkout our paper

Other design details

- How to map GaS operations to Match-Action tables
- How to select switches to form the GaS tree
- How to handle packet loss and packet out-of-order
- How to handle server or switch failure
- The practicality of SwitchTx

Implementation details

- Packet formats
- RMDA RAW_PACKET verbs for control messages
 RDMA WRITE_WITH_IMM verbs for data messages

Outline

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Experimental Setup

Hardware Platform

Server	8x Servers	
CPU	2x Intel 12-core Xeon E5-2650 CPUs	switch ₁ switch ₂ switch ₃ switch ₀
NIC	100Gbps Mellanox ConnectX-5	-
Switch	Barefoot Tofino Wedge 100BF-32X (bf-sde-8.8.1)	
		- 4 independent virtual switches

Competitors

SwitchTx	OCC + Primary-backup replication, scalable in-network coordination
FaSST [OSDI'16]	OCC + Primary-backup replication
Eris [SOSP'17]	Independent deterministic transaction, centralized in-network sequencer

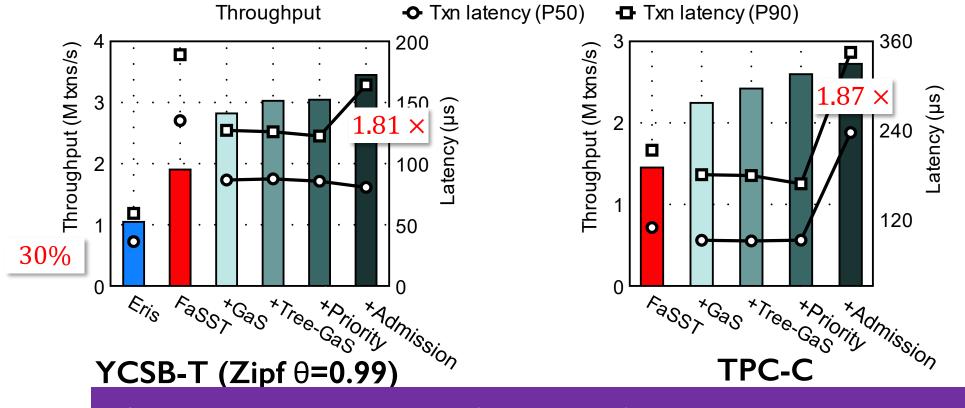
Others: Aria [VLDB'20], Calvin [SIGMOD'12] (check our paper)

Benchmarks: TPC-C, YCSB-T

Overall Performance

8 nodes, 24 threads per node

YCSB-T: a transaction reads/writes (50%:50%) 8 records, each record has an 8-byte key and a 16-byte value TPC-C: 50% New-Order + 50% Payment

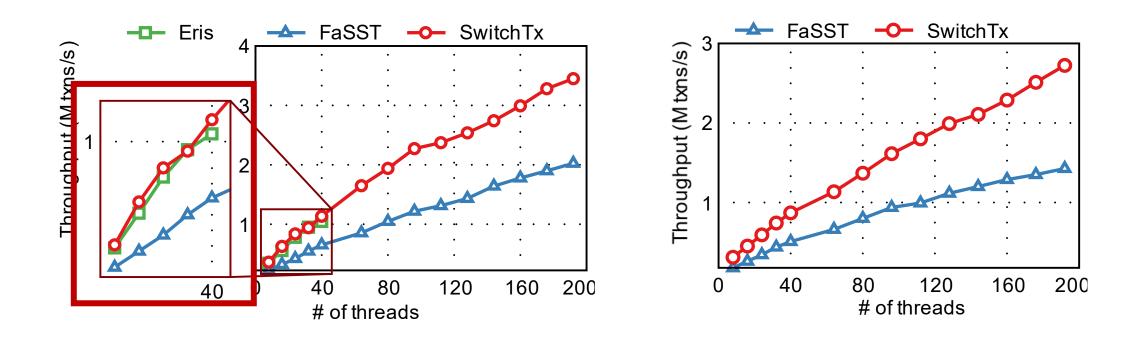


SwitchTx can boost the performance of distributed transactions

Scalability

8 nodes

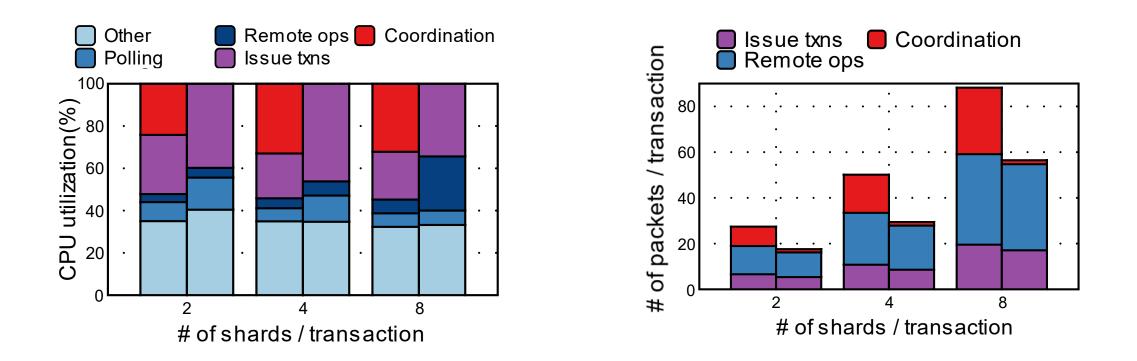
I~24 threads per node



In-switch Gather-and-Scatter is scalable

Saved CPU Resources and Packets

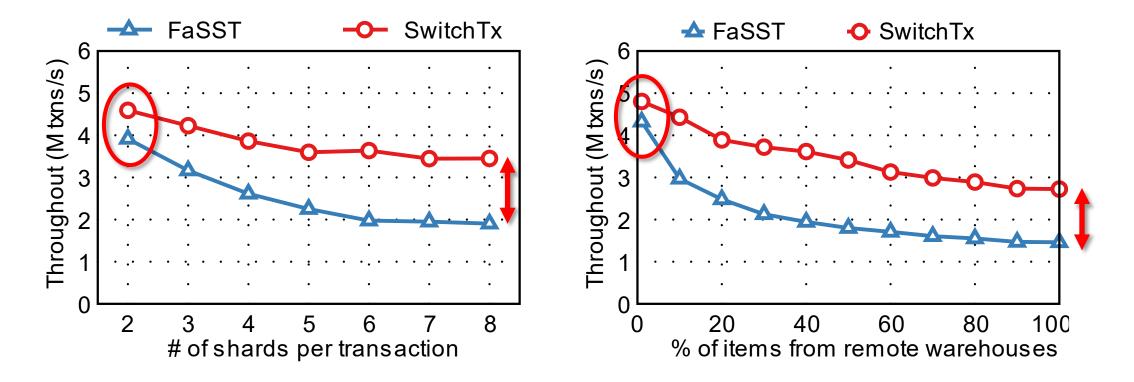
YCSB-T Benchmark



SwitchTx effectively saves CPU resources and reduces network traffic

Limitation

YCSB-T: varying the number of shards accessed by each transaction TPC-C: varying the % of remote items for New-Order transaction



SwitchTx is suitable for the transactions cross many shards

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* Goal

Reduce coordination cost in distributed transaction processing systems

* Key Idea

Using programmable switches to offload coordination tasks and manipulate transaction traffic intelligently

Techniques in SwitchTx

- Scalable in-network Gather-and-Scatter
- Priority control and dynamic admission control

Results

- SwitchTx outperforms state-of-the-arts
- SwitchTx is scalable to multiple switches



Thanks

SwitchTx: Scalable In-Network Coordination for Distributed Transaction Processing





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